Rainbow saturation and graph capacities

Dániel Korándi *

Abstract

The t-colored rainbow saturation number $\mathrm{rsat}_t(n, F)$ is the minimum size of a t-edge-colored graph on n vertices that contains no rainbow copy of F, but the addition of any missing edge in any color creates such a rainbow copy. Barrus, Ferrara, Vandenbussche and Wenger conjectured that $\mathrm{rsat}_t(n, K_s) = \Theta(n \log n)$ for every $s \geq 3$ and $t \geq \binom{s}{2}$. In this short note we prove the conjecture in a strong sense, asymptotically determining the rainbow saturation number for triangles. Our lower bound is probabilistic in spirit, the upper bound is based on the Shannon capacity of a certain family of cliques.

1 Introduction

A graph G is called *F-saturated* if it is a maximal F-free graph. The classic saturation problem, first studied by Zykov [14] and Erdős, Hajnal and Moon [4], asks for the minimum number of edges in an F-saturated graph (as opposed to the Turán problem, which asks for the maximum number of edges in such a graph). A rainbow analog of this problem was recently introduced by Barrus, Ferrara, Vandenbussche and Wenger [1], where a t-edge-colored graph is defined to be *rainbow F-saturated* if it contains no rainbow copy of F (i.e., a copy of F where all edges have different colors), but the addition of any missing edge in any color creates such a rainbow copy. Then the t-colored rainbow saturation number $\mathrm{rsat}_t(n, F)$ is the minimum size of a t-edge-colored rainbow F-saturated graph.

Among other results, Barrus et al. showed that $\Omega\left(\frac{n \log n}{\log \log n}\right) \leq \mathrm{rsat}_t(n, K_s) \leq O(n \log n)$ and conjectured that their upper bound is of the right order of magnitude:

Conjecture 1.1 ([1]). *For $s \geq 3$ and $t \geq \binom{s}{2}$, $\mathrm{rsat}_t(n, K_s) = \Theta(n \log n)$.*

Here we prove this conjecture in a strong sense: we give a lower bound that is asymptotically tight for triangles.

Theorem 1.2. *For $s \geq 3$ and $t \geq \binom{s}{2}$, we have*

$$\mathrm{rsat}_t(n, K_s) \geq \frac{t(1 + o(1))}{(t - s + 2) \log(t - s + 2)} n \log n$$

with equality for $s = 3$.

We should point out that Conjecture 1.1 was independently verified by Girão, Lewis and Popielarz [9] and by Ferrara et al. [5], but with somewhat weaker bounds. In fact, our result proves a conjecture in [9], establishing the stronger estimate $\mathrm{rsat}_t(n, K_s) = \Theta_s(\frac{n \log n}{\log t})$ with their upper bound.

*Institute of Mathematics, EPFL, Lausanne, Switzerland. Research supported in part by SNSF grants 200020-162884 and 200021-175977. Email: daniel.korandi@epfl.ch.

Our lower bound is probabilistic in spirit, using ideas of Katona and Szemerédi [10], and Füredi, Horak, Pareek and Zhu [6] (similar techniques were used in [12, 2, 11]). The upper bound for $s = 3$ is based on the following theorem that follows from a strong information-theoretic result of Gargano, Körner and Vaccaro [8] on the Shannon capacities of graph families.

Theorem 1.3. *For every $t \geq 3$, there is a set $X \subseteq [t]^k$ of $m = (t-1)^{(\frac{t-1}{t} - o(1))k}$ strings of length k from alphabet $[t] = \{1, \ldots, t\}$ such that for any $x, x' \in X$ and any $a \in [t]$, there is a position i where $x(i) \neq x'(i)$ and $x(i), x'(i) \neq a$.*

In the next section we derive Theorem 1.3 from results about the Shannon capacity of graph families. This is followed by the proof of Theorem 1.2 in Section 3.

2 Graph capacities

Let $\mathcal{G} = \{G_1, \ldots, G_r\}$ be a family of graphs on vertex set $[t]$. Let N_k be the maximum size of a set $X \subseteq [t]^k$ of strings of length k on alphabet $[t]$ such that for any two strings $x, x' \in X$ and any $G_j \in \mathcal{G}$, there is a position $i_j \in [k]$ such that $x(i_j)x'(i_j)$ is an edge in G_j. The *Shannon capacity* of the family $\mathcal{G}$ is defined as $C(\mathcal{G}) = \limsup_{k\to\infty} \frac{1}{k} \log N_k$ (see, e.g., [13, 3][1]). When $\mathcal{G} = \{G\}$, we simply write $C(G)$ for $C(\mathcal{G})$.

We need an analogous definition for strings where the occurrences of each $a \in [t]$ are proportional to some probability measure P on $[t]$. So let $\mathcal{T}^k(P, \varepsilon)$ be the set of all strings $x \in [t]^k$ such that $|\frac{1}{k}\#\{i : x(i) = a\} - P(a)| < \varepsilon$ for every $a \in [t]$, and let $M_{k,\varepsilon}$ be the maximum size of a set $X \subseteq \mathcal{T}^k(P, \varepsilon)$ such that for every $x, x' \in X$ there is an i with $x(i)x'(i) \in G$. The Shannon capacity within type P is $C(G, P) = \lim_{\varepsilon\to 0} \limsup_{k\to\infty} \frac{1}{k} \log M_{k,\varepsilon}$. Using a clever construction, Gargano, Körner and Vaccaro [8] showed that $C(\mathcal{G})$ can be expressed in terms of the $C(G_j, P)$:

Theorem 2.1 ([8])**.** *For a family of graphs $\mathcal{G} = \{G_1, \ldots, G_r\}$ on vertex set $[t]$, we have*

$$C(\mathcal{G}) = \max_P \min_{G_j \in \mathcal{G}} C(G_j, P).$$

In fact, they proved a more general result for *Sperner capacities*, the analogous notion for directed graphs. What we need is a corollary that follows easily from this theorem using standard tools about graph entropy (see the survey of Simonyi [13] for more information). Here we give a self-contained argument that goes along the lines of a proof by Gargano, Körner and Vaccaro [7] of the case $s = 2$.

Corollary 2.2. *Let $2 \leq s \leq t$ be an integer and let $\mathcal{G}$ be the family of all s-cliques on $[t]$ (each with $t - s$ isolated vertices). Then $C(\mathcal{G}) = \frac{s}{t} \log s$.*

Proof. For the lower bound, we can take P to be the uniform measure on $[t]$. Then by Theorem 2.1, it is enough to show that $C(G, P) \geq \frac{s}{t} \log s$ where G is a clique on $[s]$ with isolated vertices $s+1, \ldots, t$. Let $X_k \subseteq \mathcal{T}^k(P, \frac{1}{k})$ be the set of all strings x of length k such that the first $\lfloor sk/t \rfloor$

[1]The usual definition is with binary logarithm, but the base of our logarithms is unimportant for our purposes.

letters of x contain $\lfloor k/t \rfloor$ or $\lceil k/t \rceil$ instances of each $a \in [s]$, and $x(i) = b$ for every $s+1 \le b \le t$ and $\frac{(b-1)k}{t} < i \le \frac{bk}{t}$. Then

$$C(G,P) \ge \lim_{k\to\infty} \frac{\log(X_k)}{k} = \lim_{k\to\infty} \frac{1}{k} \log \frac{(\frac{sk}{t})!}{((\frac{k}{t})!)^s} = \lim_{k\to\infty} \frac{1}{k} \log(s^{sk/t}) = \frac{s}{t} \log s.$$

For the upper bound, let $X \subseteq [t]^k$ be a maximum set of strings such that for any $x, x' \in X$ and for every s-clique $G \in \mathcal{G}$, there is an $i \in [k]$ such that $x(i)x'(i) \in G$. We set $m = |X|$ to be this maximum. We may assume that $\{1, \ldots, s\}$ are the s least frequent elements appearing in the strings of X. Let d_x be the number of elements in $x \in X$ that are not in $[s]$, so $\sum_{x\in X} d_x \ge \frac{t-s}{t} mk$, and let X_x be the set of strings obtained from x by replacing these elements arbitrarily with numbers from $[s]$. Then $|X_x| = s^{d_x}$, and $X_x, X_{x'}$ are disjoint for distinct $x, x' \in X$ because any string from X_x will differ from any string in $X_{x'}$ at the position i where $x(i)x'(i)$ is an edge of the clique on $[s]$. Then using Jensen's inequality we have

$$s^k \ge \sum_{x\in X} s^{d_x} \ge m \cdot s^{(\sum_{x\in X} d_x)/m} \ge m \cdot s^{\frac{(t-s)k}{t}},$$

and hence $m \le s^{sk/t}$, implying $C(\mathcal{G}) \le \frac{1}{k} \log m \le \frac{s}{t} \log s$. □

Theorem 1.3 clearly follows from the case $s = t - 1$.

3 Rainbow saturation

Proof of Theorem 1.2. For the lower bound, suppose H is a t-edge-colored rainbow K_s-saturated graph, and split its vertices into two parts: let $A = \{a_1, \ldots, a_k\}$ be the set of vertices of degree at least $d = \log^3 n$, and B be the rest. We may assume $|A| \le \frac{n}{\log n}$ (otherwise H has at least $\frac{1}{2} n \log^2 n$ edges), and thus B contains $m \ge (1 - \frac{1}{\log n})n$ vertices. Now let us define a string $x_v \subseteq [t+1]^k$ for every $v \in B$ that encodes the colors of the A-B edges touching v as follows: $x_v(i)$ is $t+1$ if $a_i v$ is not an edge in H, otherwise it is the color of $a_i v$.

Assume, without loss of generality, that $t-s+3, \ldots, t$ are the $s-2$ most common colors among the A-B edges. For $v \in B$, let $X_v \subseteq [t-s+2]^k$ be the set of strings obtained from x_v by replacing each $t-s+3, \ldots, t+1$ with an arbitrary number from $[t-s+2]$. Then if d_v denotes the number of A-B edges in H touching v and d'_v denotes the number of such edges of colors $t-s+3, \ldots, t$, then $|X_v| = (t-s+2)^{k-d_v+d'_v}$.

We claim that if $v, w \in B$ are non-adjacent with no common neighbor in B, then X_v and X_w have no string in common. Indeed, adding the edge vw of color t creates a rainbow K_s with $s-2$ vertices in A. So there must be an a_i such that $a_i v$ and $a_i w$ have different colors, also differing from $t-s+3, \ldots, t$. But then all the strings in X_v have the color of $a_i v$ as their i'th letter, and all the strings in X_w have the color of $a_i w$ as their i'th letter, so X_v and X_w are disjoint.

Since vertices in B have degree at most d, each $v \in B$ has at most d^2 vertices $w \in B$ that are either adjacent to v or have a common neighbor with v in B. So each string in $[t-s+2]^k$ can appear in no more than d^2+1 collections X_w, and hence we get

$$(d^2+1)(t-s+2)^k \geq \sum_{v\in B} |X_v| = \sum_{v\in B}(t-s+2)^{k-d_v+d'_v}$$
$$d^2+1 \geq \sum_{v\in B}(t-s+2)^{d'_v-d_v} \geq m\cdot(t-s+2)^{\frac{1}{m}(\sum_{v\in B} d'_v - \sum_{v\in B} d_v)}$$

using Jensen's inequality.

Now $t-s+3,\dots,t$ were the $s-2$ most common colors, so we also have $\sum_{v\in B} d'_v \geq \frac{s-2}{t}\sum_{v\in B} d_v$ and thus $\sum_{v\in B} d'_v - \sum_{v\in B} d_v \geq \frac{s-2-t}{t}\sum_{v\in B} d_v$. Taking logs, we obtain

$$\sum_{v\in B} d_v \geq \frac{t}{t-s+2} m\left(\log_{t-s+2} m - \log_{t-s+2}(d^2+1)\right).$$

As the left-hand side is a lower bound on the number of edges in H, this establishes the desired lower bound (using $d=\log^3 n$ and $m=n+o(n)$).

For the upper bound in the case of triangles, let k be large enough, and take a set X of size m as provided by Theorem 1.3. Consider a k-by-m complete bipartite graph G_0 with parts A and B, where $A=\{a_1,\dots,a_k\}$, and B corresponds to the strings in X. For every vertex $v\in B$, we look at the corresponding string $x\in X$, and color each edge va_i by the color $x(i)$. G_0 is clearly (rainbow) triangle-free, and by the definition of X, adding an edge to G_0 between two vertices of B in any color $a\in[t]$ creates a rainbow triangle.

Now let G be a maximal rainbow triangle-free supergraph of G_0. Then G is rainbow triangle-saturated by definition, and compared to G_0, it only has new edges induced by A, thus it has at most $km+\binom{k}{2}$ edges. Here $n=k+m$ and $k=\frac{t(1+o(1))}{(t-1)\log(t-1)}\log m$, implying the required upper bound. □

For $s>3$ our lower bound is probably not tight. It would be interesting to determine the asymptotics of $\mathrm{rsat}_t(n,K_s)$ for general s.

Acknowledgements. I thank Shagnik Das for finding [7] for me, and Gábor Simonyi for some clarifications about capacities.

References

[1] M. D. Barrus, M. Ferrara, J. Vandenbussche and P. S. Wenger, Colored saturation parameters for rainbow subgraphs, *J. Graph Theory*, **86** (2017), 375-386.

[2] B. Bollobás and A. Scott, Separating systems and oriented graphs of diameter two, *J. Combin. Theory Ser. B* **97** (2007), 193-203.

[3] I. Csiszár and J. Körner, Information Theory, 2nd edition, Cambridge University Press, 2011.

[4] P. Erdős, A. Hajnal and J.W. Moon, A problem in graph theory, *Amer. Math. Monthly*, **71** (1964), 1107-1110.

[5] M. Ferrara, D. Johnston, S. Loeb, F. Pfender, A. Schulte, H. C. Smith, E. Sullivan, M. Tait and C. Tompkins, On edge-colored saturation problems, arXiv:1712.00163 preprint

[6] Z. Füredi, P. Horak, C. M. Pareek and X. Zhu, Minimal oriented graphs of diameter 2, *Graphs Combin.* **14** (1998), 345-350.

[7] L. Gargano, J. Körner and U. Vaccaro, Sperner capacities, *Graphs Combin.*, **9** (1993), 31-46.

[8] L. Gargano, J. Körner and U. Vaccaro, Capacities: from information theory to extremal set theory, *J. Combin. Theory Ser. A*, **68** (1994), 296-316.

[9] A. Girão, D. Lewis and K. Popielarz, Rainbow saturation of graphs, arXiv:1710.08025 preprint

[10] G. Katona and E. Szemerédi, On a problem of graph theory, *Studia Sci. Math. Hungar.* **2** (1967), 23-28.

[11] D. Korándi and B. Sudakov, Saturation in random graphs, *Random Structures Algorithms* **51** (2017), 169-181.

[12] A. V. Kostochka, T. Łuczak, G. Simonyi and E. Sopena, On the minimum number of edges giving maximum oriented chromatic number, in: Contemporary Trends in Discrete Mathematics, *DIMACS Series in Discrete Mathematics and Theoretical Computer Science*, vol 49. (1999), 179-182.

[13] G. Simonyi, Perfect graphs and graph entropy. An updated survey, in: Perfect Graphs, Wiley (2001), 293-328.

[14] A. Zykov, On some properties of linear complexes (in Russian), *Mat. Sbornik N. S.* **24** (1949), 163-188.