

On Superspherical Graphs

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ABSTRACT

In this paper we consider two conjectures of H. M. Mulder [2] and give partial solutions for them. First we prove that a triangle free superspherical graph is interval regular provided it satisfies an additional (not too strong) condition. Furthermore, we show that a spherical interval regular graph is interval monotone.

1. Introduction

Graphs that are well-structured or highly symmetrical are studied extensively. In particular hypercubes and graphs that are close to hypercubes draw much attention. In his book [2] H.M. Mulder put forward the following conjecture (see the definitions below).

Conjecture 1.1. *An interval regular graph is interval monotone.*

Furthermore, he proposed the following.

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Conjecture 1.2. *A triangle free superspherical graph is interval regular.*

In the present paper we prove two related theorems. In order to formulate the results we need some definitions.

Definition 1.3. *Let $G = (V, E)$ be a finite graph and let $u, v \in V$. The interval $I(u, v)$ is the subgraph of G induced by the set of all vertices lying on a shortest path between u and v . The length of the interval is the distance $d(u, v)$ of u and v .*

Let us denote the d -dimensional hypercube by \mathcal{B}_d . So \mathcal{B}_d is the graph whose vertex set is $\{0, 1\}^d$ and two vertices are connected by an edge iff they differ in exactly one position.

Definition 1.4. *Let $G = (V, E)$ be a finite connected graph. G is called interval regular if for any $u, v \in V$ the subgraph induced by the set of edges between levels in the interval $I(u, v)$ is a hypercube \mathcal{B}_d with $d = d(u, v)$.*

Note that this definition is different from that of [2], but the two are equivalent. Furthermore, $I(u, v)$ need not be isomorphic to \mathcal{B}_d even in the case of an interval regular G . Let us denote the property that *the subgraph induced by the set of edges between levels in the interval $I(u, v)$ is a hypercube \mathcal{B}_d with $d = d(u, v)$* by $I(u, v) \asymp \mathcal{B}_d$. The next definition deals with convexity properties of intervals.

Definition 1.5. *Let $G = (V, E)$ be a finite graph and let $A \subset V$. We say that A is convex if for every $x, y \in A$ the vertices of $I(x, y)$ are contained in A . G is interval monotone if each interval of G is convex. G is said to have the quadrangle property, if for every interval $I(u, v)$ of G and every $x, y \in I(u, v)$ such that $d(x, u) = 1$ and $d(y, u) = 1$, there exists a $u \neq z \in I(u, v)$ such that $d(z, x) = d(z, y) = 1$.*

One more definition is needed.

Definition 1.6. *Let $G = (V, E)$ be a finite graph and let $u, v \in V$ such that $d(u, v) = d$. Let $x, y \in I(u, v)$. We say that x and y are diametrical if $d(x, y) = d(u, v) = d$. G is called spherical if for every point in each of its intervals there exists at least one diametrical point in that particular interval. G is superspherical if spherical and for every point the diametrical pair is unique in every interval.*

If x and y are diametrical in some interval, then y is often called the opposite or complement of x in that interval. It is easy to see that \mathcal{B}_d is

an example of superspherical graphs. However, a superspherical graph need not be a hypercube as the following example shows. Let $\mathcal{B}_d(1, d)$ denote the graph whose vertex set is $\{0, 1\}^d$ and two vertices are connected by an edge iff they differ in exactly one or d positions. Then $\mathcal{B}_{2t}(1, 2t)$ is a superspherical graph. Figure 1 depicts $\mathcal{B}_4(1, 4)$.

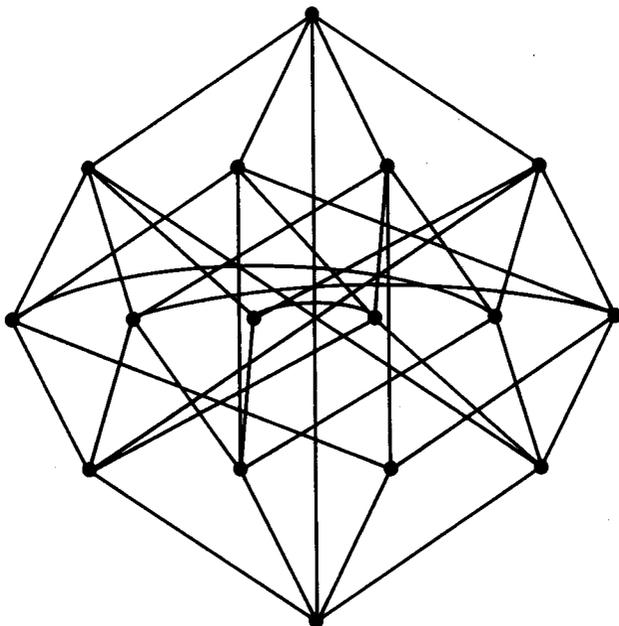


Fig. 1. $\mathcal{B}_4(1, 4)$

Now we can formulate the main results.

Theorem 1.7. *If a triangle free superspherical graph G satisfies the quadrangle property, then it is interval regular.*

Theorem 1.8. *If an interval regular graph G is spherical, then it is interval monotone.*

As an immediate corollary we obtain the following.

Corollary 1.9. *For triangle free superspherical graphs the following three conditions are equivalent:*

- (i) G satisfies the quadrangle property
- (ii) G is interval monotone
- (iii) G is interval regular. ■

We mention that Havel and Liebl proved that a bipartite superspherical graph is in fact isomorphic to a hypercube [1].

2. Proofs

The proof of Theorem 1.7 consists of several small steps. We use induction on d . Let α be a neighbour of v in $I(u, v)$. Let $\bar{\alpha}$ be its unique diametrical pair. Let level i consist of vertices in $I(u, v)$ of distance i from u . For $z \in I(u, v)$ let $l(z)$ denote the level of z . The first step is the case $d = 2$.

Step 1. *If x and y are such that $d(x, y) = 2$, then $I(x, y) \cong \mathcal{B}_2$.*

Proof. $I(x, y)$ consists of common neighbours of x and y besides x and y because $d(x, y) = 2$. By the triangle free property there does not exist edge between two different neighbours of x . There is at least one vertex $z \in I(x, y)$, so there must exist its diametrical pair, say w . Now, if there were any other vertex in $I(x, y)$, then z would have more than one diametrical pair, a contradiction. ■

Step 2. $I(u, \alpha) \cap I(\bar{\alpha}, v) = \emptyset$.

Proof. Suppose in contrary that there exists β of level i in $I(u, v)$ such that $\beta \in I(u, \alpha) \cap I(\bar{\alpha}, v)$. Then we have $d(\bar{\alpha}, \beta) = i - 1$ and $d(\beta, \alpha) = d - i$ so

$$d(\alpha, \bar{\alpha}) \leq d(\bar{\alpha}, \beta) + d(\beta, \alpha) = i - 1 + d - i < d$$

would hold, a contradiction. ■

Step 3. *There does not exist an $x \in (I(u, v) \setminus (I(u, \alpha) \cup I(\bar{\alpha}, v)))$ of level i such that it has a neighbour of level $i - 1$ in $I(u, \alpha)$, and one of level $i + 1$ in $I(\bar{\alpha}, v)$.*

Proof. Suppose that there exists such an x and let its neighbours be z and w ($l(z) = i - 1$ and $l(w) = i + 1$). Then $d(z, w) = 2$, so $I(z, w) \cong \mathcal{B}_2$. Hence, there exists a vertex t such that $t \in I(z, w) \setminus \{x\}$. It is easy to see that t is of level i in $I(u, v)$. Suppose first, that $w \neq v$ and $z \neq u$. If $t \notin I(\bar{\alpha}, v)$, then w would have at least $i + 2$ neighbours of level i in $I(u, v)$ that contradicts to the induction hypothesis $I(u, w) \bowtie \mathcal{B}_{i+1}$. On the other hand, if $t \in I(\bar{\alpha}, v)$, then z has at least $d - i + 2$ neighbours on level i . If $z \neq u$, then this

contradicts to the induction hypothesis. Now let $w = v$. Pick a neighbour q of v in $I(\bar{\alpha}, v)$. Let a be a common neighbour of x and q on level $d - 2$, and let b be that of α and q provided by the quadrangle property. Then $q \notin I(u, \alpha) \cup I(\bar{\alpha}, v)$ and $b \in I(u, \alpha)$. So q has at least d neighbours on level $d - 2$ in $I(u, v)$ that contradicts to the induction hypothesis $I(u, q) \bowtie \mathcal{B}_{d-1}$. For $z = u$ we obtain contradiction in similar way as in the case $w = v$. ■

Step 4. *If $l(y) = j \neq d$ and all neighbours of y of level $j - 1$ are in $I(u, v) \setminus I(u, \alpha)$, then $I(u, \alpha) \cap I(u, y) = \{u\}$.*

Proof. Suppose that $z \neq u$, such that $z \in I(u, \alpha) \cap I(u, y)$ and z is of the highest level subject to this condition. Let $l(z) = k$. Then z has $d - 1 - k$ neighbours of level $k + 1$ in $I(u, \alpha)$. Furthermore, z has $j - k$ neighbours of level $k + 1$ in $I(u, y)$. All these are different, so z has at least $d + j - 2k - 1 > d - k$ neighbours of level $k + 1$ in $I(u, v)$ that contradicts to the induction hypothesis $I(z, v) \bowtie \mathcal{B}_{d-k}$. ■

Note 1. *Similar proposition holds if we consider level $j + 1$ neighbours with the conclusion that $I(\bar{\alpha}, v) \cap I(y, v) = \{v\}$.*

Step 5. *If $z \in I(u, \alpha)$, $z \neq u$, $l(z) = i$, $w \in I(u, v) \setminus I(u, \alpha)$, $l(w) = i + 1$ and $\{z, w\}$ is an edge, then $w \in I(\bar{\alpha}, v)$. The same holds for changing $I(u, \alpha)$ to $I(\bar{\alpha}, v)$ and $i + 1$ to $i - 1$.*

Proof. Suppose that $w \notin I(\bar{\alpha}, v)$. Then no neighbour of w of level $i + 2$ can be in $I(\bar{\alpha}, v)$ by Step 3. Thus, by Step 4 we have that $I(\bar{\alpha}, v) \cap I(w, v) = \{v\}$. Let $w = t_0, t_1, \dots, t_s = v$ be a shortest $(d - i - 1)$ -long path from w to v . Because $I(u, t_{s-1}) \cap I(u, \alpha) \ni z \neq u$, t_{s-1} must have a neighbour of level $d - 2$ in $I(u, \alpha)$. This contradicts to Step 3. ■

Note 2. *If $z \in I(u, \alpha)$ ($z \in I(\bar{\alpha}, v)$) such that $l(z) = i$, then all neighbours of z of level $i - 1$ ($i + 1$) are in $I(u, \alpha)$ ($I(\bar{\alpha}, v)$).*

Step 6. *For every $z \in I(u, \alpha)$, such that $z \neq u$ and $l(z) = i$, there exists exactly one w $l(w) = i + 1$ such that $\{z, w\}$ is an edge and $w \notin I(u, \alpha)$. The same holds for exchanging $I(u, \alpha)$ to $I(\bar{\alpha}, v)$ and $i + 1$ to $i - 1$.*

Proof. $I(z, v) \bowtie \mathcal{B}_{d-i}$ by the induction hypothesis ($z \neq u$), so z has $d - i$ neighbours of level $i + 1$ in $I(u, v)$. Out of those $d - i - 1$ lie in $I(u, \alpha)$. ■

Let us denote the above w by z' and let $u' = \bar{\alpha}$ by definition. Then by Step 3-5, $z' \in I(\bar{\alpha}, v)$ if $z \in I(u, \alpha)$ and vica versa. Furthermore, $(z')' = z$.

Hence, $z \mapsto z'$ is a bijection between $I(u, \alpha)$ and $I(\bar{\alpha}, v)$.

Step 7. $z \mapsto z'$ is an isomorphism between the graphs induced by edges between different levels of $I(u, \alpha)$ and $I(\bar{\alpha}, v)$, respectively.

Proof. It is enough to show that $\{x, y\}$ is an edge in $I(u, \alpha)$ implies $\{x', y'\}$ is an edge in $I(\bar{\alpha}, v)$. The reverse comes from $(z')' = z$. Let $\{x, y\}$ be an edge, such that $l(x) = i$ and $l(y) = i + 1$. Then $d(x, y') = d(x', y) = 2$, so x and y' must have a unique common neighbour other than y . If it is in $I(\bar{\alpha}, v)$, then by Step 6, it must be x' . If it is in $I(u, \alpha)$, then it must be on level $i + 1$. Then again by Step 6, it must be y , a contradiction. If it is in $I(u, v) \setminus (I(u, \alpha) \cup I(\bar{\alpha}, v))$, then it contradicts to Step 3. ■

Step 8. If $l(z) = i$, $z \in I(u, \alpha)$, $l(w) \leq i$ and $w \in I(\bar{\alpha}, v)$, then $\{z, w\}$ is not an edge.

Proof. By the triangle inequality we have that

$$\begin{aligned} (\alpha, \bar{\alpha}) &\leq d(\alpha, z) + d(z, w) + d(w, \bar{\alpha}) \\ &\leq d - 1 - i + 1 + i - 1 = d - 1 \end{aligned}$$

a contradiction. ■

Now, if we do not consider edges between vertices of the same level, then $I(u, \alpha)$ and $I(\bar{\alpha}, v)$ are $d - 1$ dimensional hypercubes, respectively. Furthermore, the latter one is positioned one level higher, and there is an edge between corresponding vertices of $I(u, \alpha)$ and $I(\bar{\alpha}, v)$. No other edge is going between vertices of different levels so we can conclude that $I(u, \alpha) \cup I(\bar{\alpha}, v) \bowtie \mathcal{B}_d$.

Step 9. $I(u, v) = I(u, \alpha) \cup I(\bar{\alpha}, v)$.

Proof. Suppose in contrary that there exists a neighbour x of v that is in $I(u, v) \setminus (I(u, \alpha) \cup I(\bar{\alpha}, v))$. Then by the quadrangle property there exists $z \in I(u, v)$ of level $d - 2$ such that z is connected to both x and α . This implies that $z \in I(u, \alpha)$. However, this contradicts to Step 3. ■

Proof of Theorem 1.8. Let $u, v \in V(G)$ and $x, y \in I(u, v)$. Let us denote the distance of x and y in the graph $I(u, v)$ by $d_{uv}(x, y)$ to distinguish from their distance in G , $d(x, y)$. If $d_{uv}(x, y) = d(x, y) = b$, then $I(x, y) \bowtie \mathcal{B}_b$. However, $I(u, v)$ contains a b -cube between x and y because $I(u, v) \bowtie \mathcal{B}_d$.

This means that $I(x, y) \subset I(u, v)$. On the other hand, if $d_{uv}(x, y) > d(x, y)$, then x does not have a diametrical pair in $I(u, v)$, which contradicts to the assumption that G is spherical. ■

References

- [1] I. Havel and P. Liebl, unpublished result (Prague).
- [2] H. M. Mulder, *The Interval Function of a Graph*, Math. Centre Tracts, **132** Amsterdam 1980.

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