

Improved Ramsey-type results in comparability graphs

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EPFL

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joint work with István Tomon

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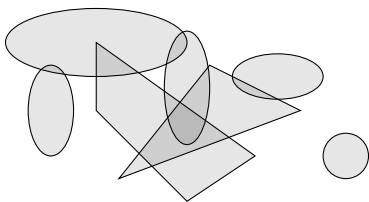
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“Dilworth”:

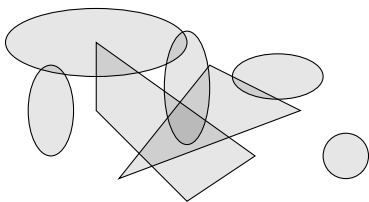
If X has no **t -chain**, then it can be partitioned into t **antichains**.

Convex sets in the plane

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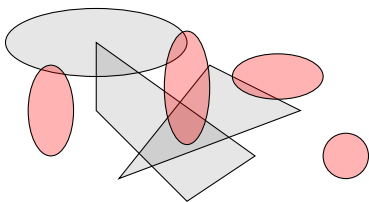
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Question

Given n convex sets in the plane, how big is the largest **disjoint** or **pairwise intersecting** subfamily?

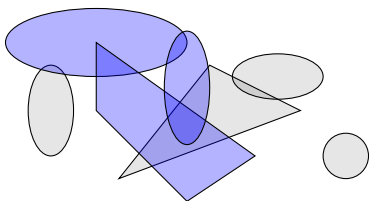
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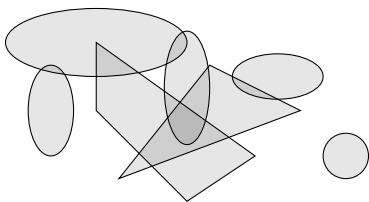
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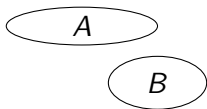
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Theorem (Larman-Matoušek-Pach-Törőcsik, 1994)

At least $n^{1/5}$.

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A, B disjoint, B is “below” A



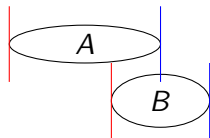
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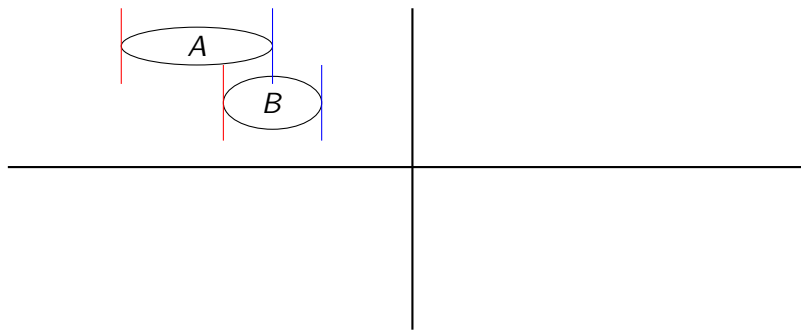
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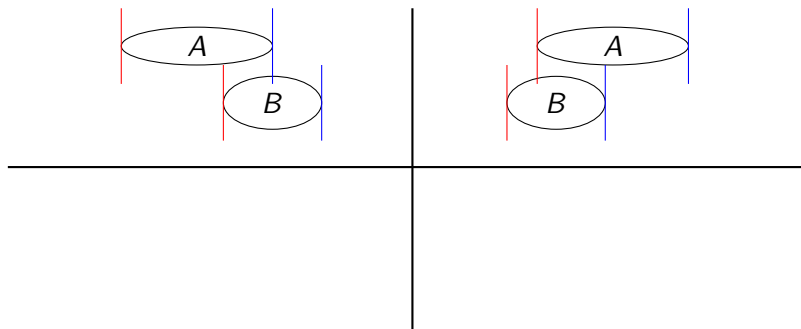
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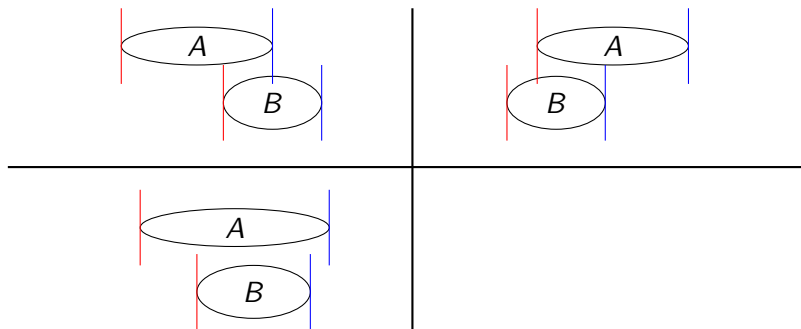
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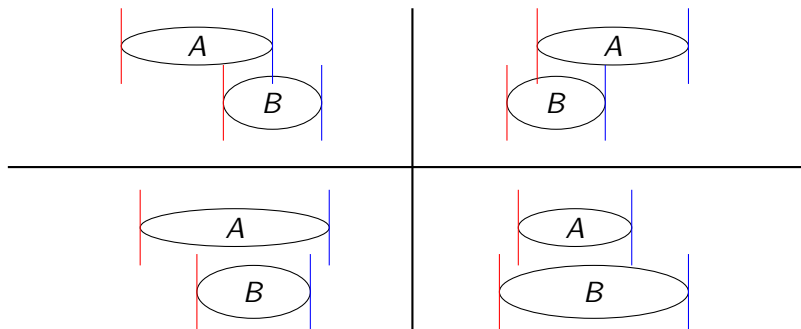
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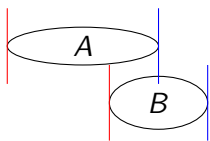
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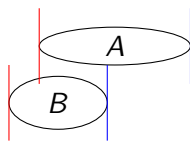
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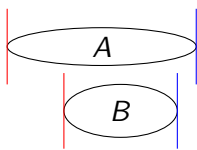
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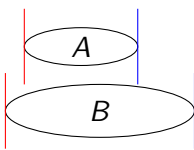
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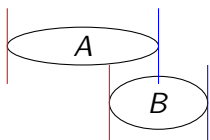
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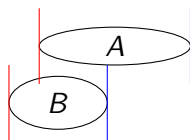
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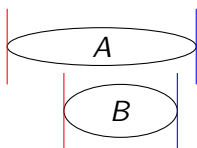
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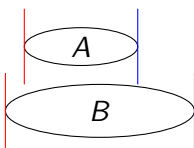
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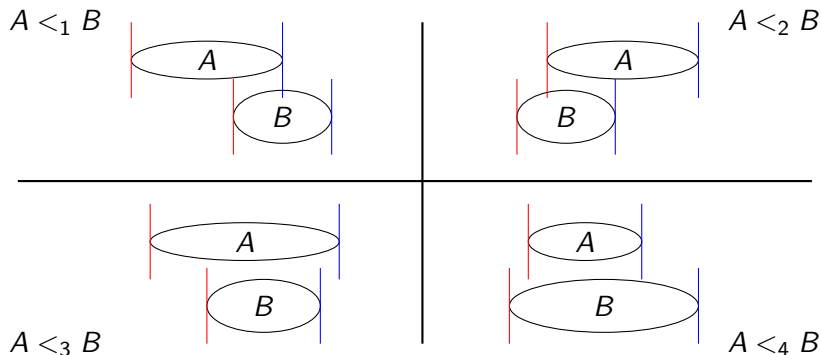
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$<_1, <_2, <_3, <_4$ are partial orders

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Sets A, B are disjoint **iff** they are comparable in any of $\langle_1, \dots, \langle_4$.

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$\Rightarrow S_4$ incomparable in all 4 posets \rightarrow intersecting family.

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Comparability graph of a poset

Connect a, b with an edge if $a < b$ or $b < a$.

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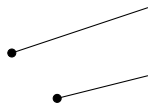
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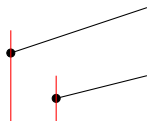
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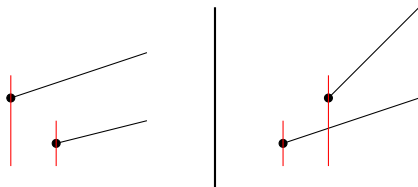
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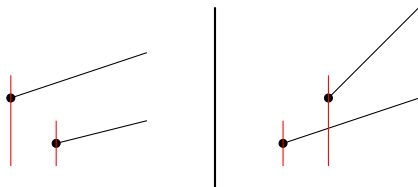
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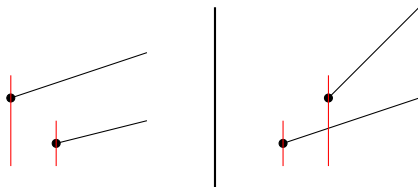
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Lower bound

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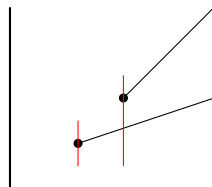
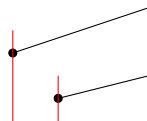
$$n^{1/5}$$

$$n^{0.405} \text{ (Kynčl)}$$

halflines

$$n^{1/3}$$

$$n^{0.431} \text{ (LMPT)}$$



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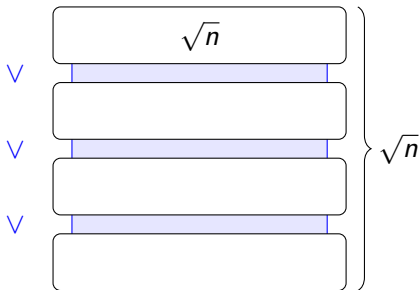
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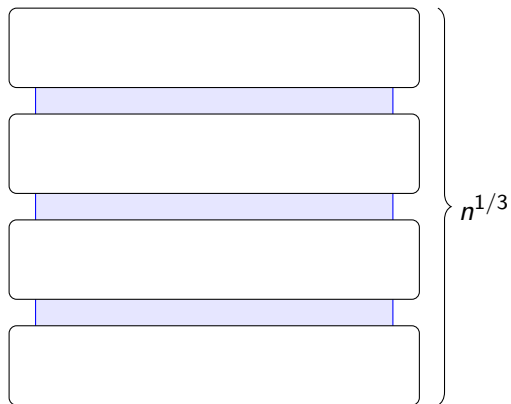
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Theorem (K-Tomon, 2019+)

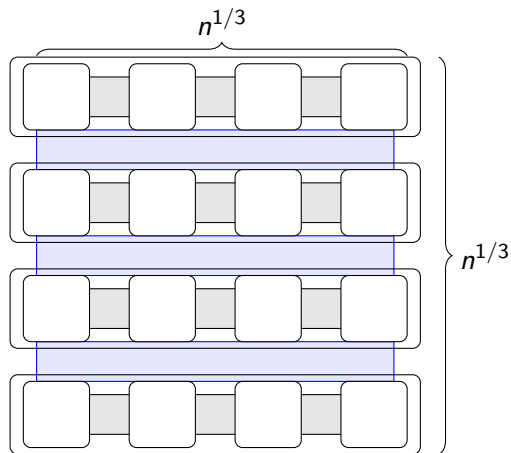
$$f_2(n) = n^{1/3+o(1)}$$

Construction for $f_2(n) \leq n^{1/3+o(1)}$

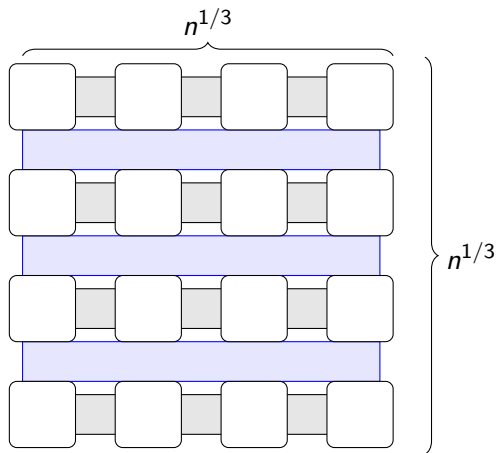
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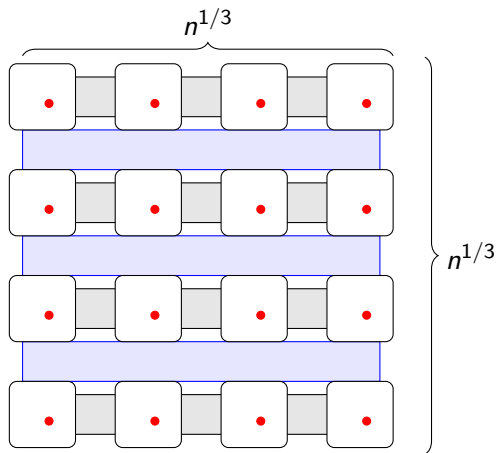
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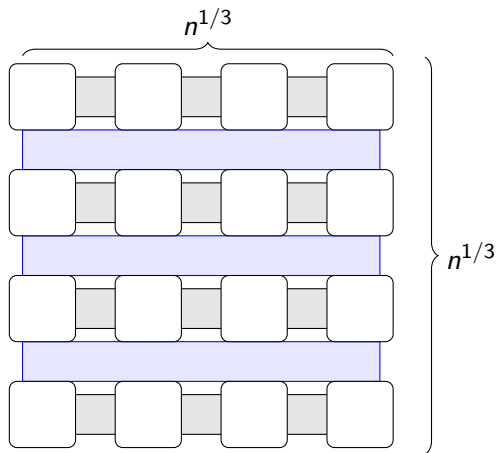
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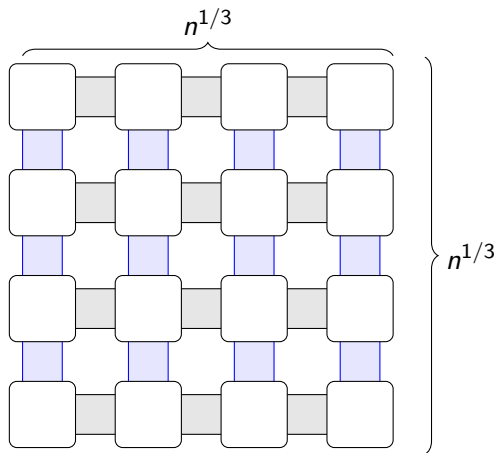
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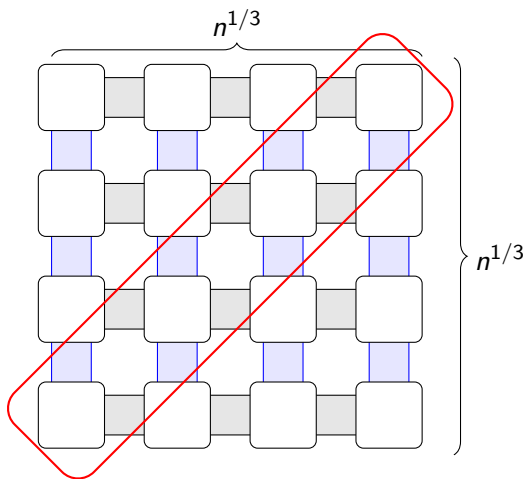
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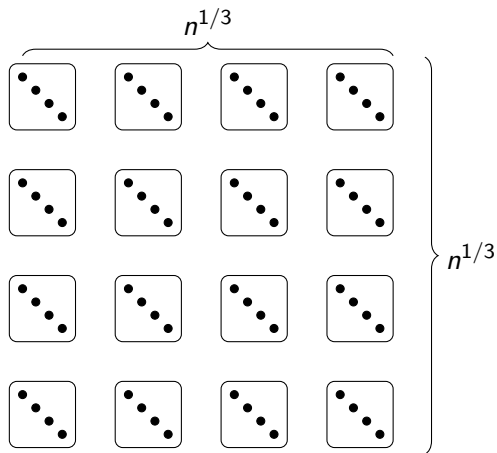
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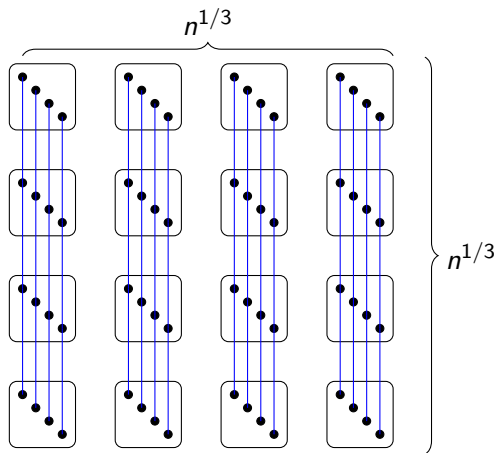
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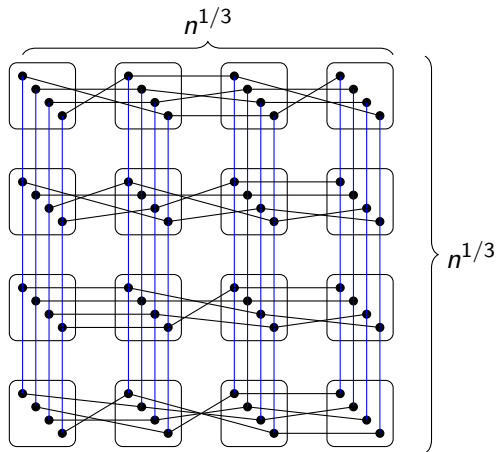
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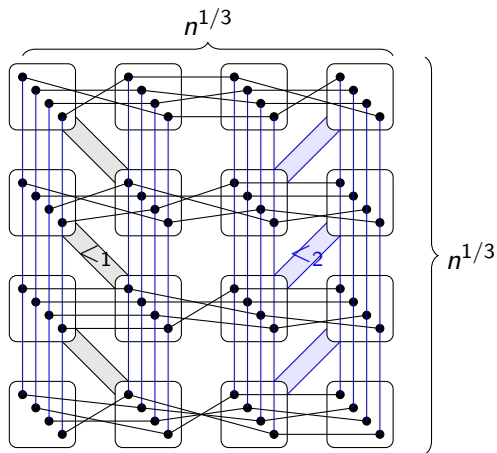
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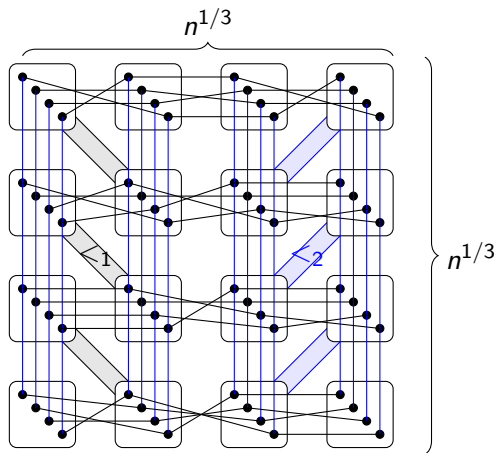
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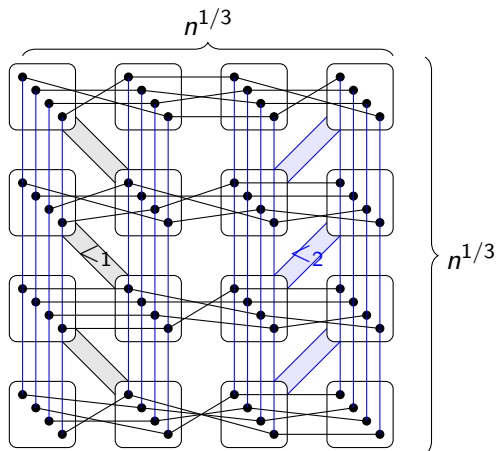
Construction for $f_2(n) \leq n^{1/3+o(1)}$

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Claim 2:

$$\omega(G) \leq n^{1/3} \cdot \frac{\log n}{\log \log n}$$



Construction for $f_2(n) \leq n^{1/3+o(1)}$

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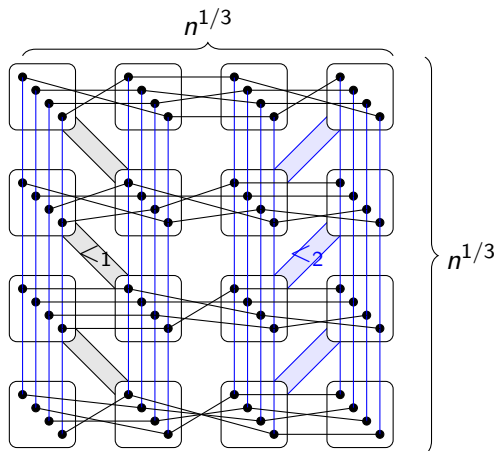
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Claim 2:

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Claim 1:

- every independent set is within a row or column
- hence it is covered by $n^{1/3}$ chains



Construction for $f_2(n) \leq n^{1/3+o(1)}$

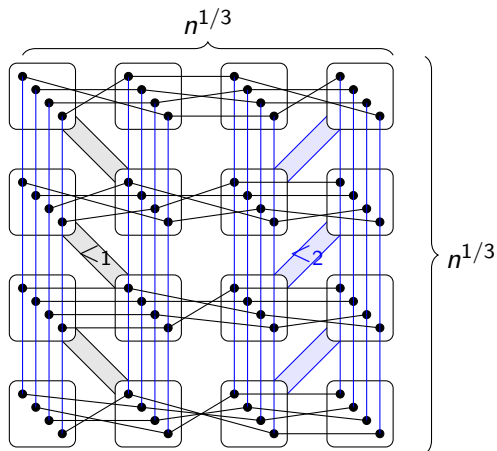
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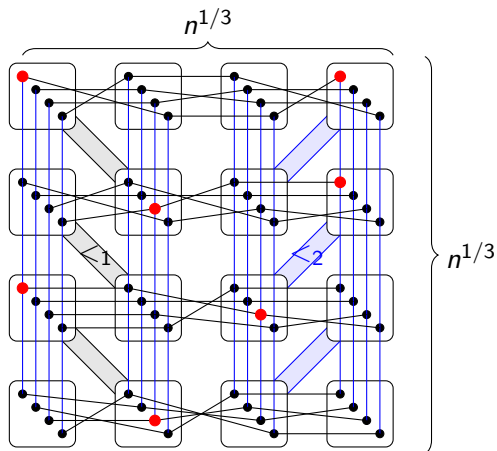
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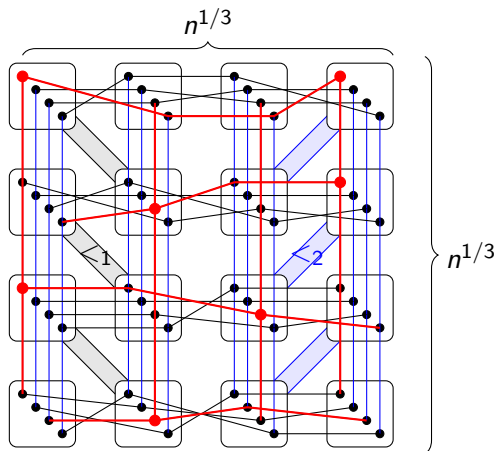
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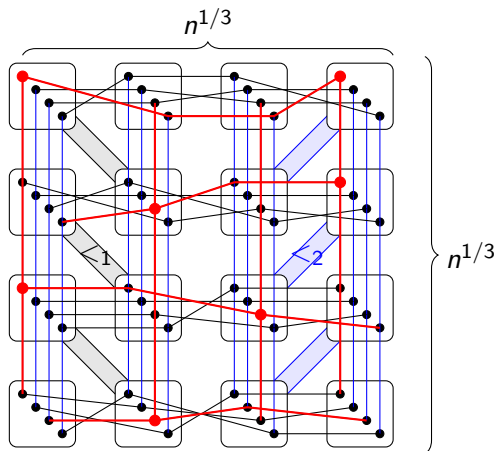
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- every clique corresponds to a fixed chain in each row and column

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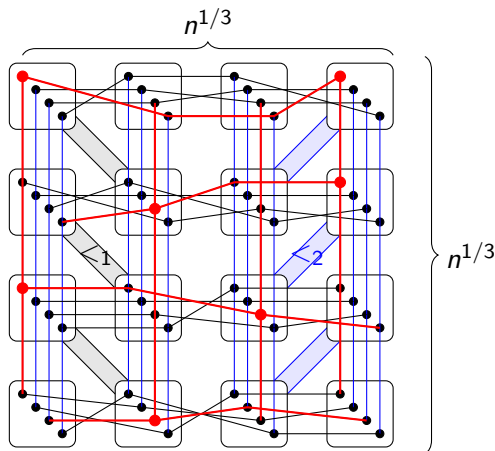
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Claim 2:

$$\omega(G) \leq n^{1/3} \cdot \frac{\log n}{\log \log n}$$

Claim 2:

- every clique corresponds to a fixed chain in each row and column
- $\mathbf{P}(\text{fixed row chain intersects fixed column chain}) = n^{-1/3}$



Remarks

- ▶ Improving $n^{1/3}$ for halflines needs different approach

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Theorem (K-Tomon, 2019+)

$$g_k(n) \leq \frac{n}{(\log n)^k}$$

Thank you!